

# Distributed Systems

## 17. MapReduce

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# Credit

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Much of this information is from the Google Code University:

<http://code.google.com/edu/parallel/mapreduce-tutorial.html>

See also:

<http://hadoop.apache.org/common/docs/current/>

for the Apache Hadoop version

Read this (the definitive paper):

<http://labs.google.com/papers/mapreduce.html>

# Background

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- Traditional programming is serial
- Parallel programming
  - Break processing into parts that can be executed concurrently on multiple processors
- Challenge
  - Identify tasks that can run concurrently and/or groups of data that can be processed concurrently
  - Not all problems can be parallelized

# Simplest environment for parallel processing

- No dependency among data
- Data can be split into equal-size chunks
- Each process can work on a chunk
- Master/worker approach
  - Master:
    - Initializes array and splits it according to # of workers
    - Sends each worker the sub-array
    - Receives the results from each worker
  - Worker:
    - Receives a sub-array from master
    - Performs processing
    - Sends results to master

# MapReduce

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- Created by Google in 2004
  - Jeffrey Dean and Sanjay Ghemawat
- Inspired by LISP
  - Map(function, set of values)
    - Applies function to each value in the set  
`(map 'length '(()) (a) (a b) (a b c))) ⇒ (0 1 2 3)`
  - Reduce(function, set of values)
    - Combines all the values using a binary function (e.g., +)  
`(reduce #' + '(1 2 3 4 5)) ⇒ 15`

# MapReduce

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- MapReduce
  - Framework for parallel computing
  - Programmers get simple API
  - Don't have to worry about handling
    - parallelization
    - data distribution
    - load balancing
    - fault tolerance
- Allows one to process huge amounts of data (terabytes and petabytes) on thousands of processors

# Who has it?

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- Google:
  - Original proprietary implementation
- Apache Hadoop MapReduce
  - Most common (open-source) implementation
  - Built to specs defined by google
- Amazon Elastic MapReduce
  - Uses Hadoop MapReduce running on Amazon EC2

# MapReduce

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- **Map**: (input shard)  $\rightarrow$  intermediate(key/value pairs)
  - Map calls are distributed across machines by automatically partitioning the input data into M "shards".
  - MapReduce library groups together all intermediate values associated with the same intermediate key & passes them to the *Reduce* function
- **Reduce**: intermediate(key/value pairs)  $\rightarrow$  result files
  - Accepts an intermediate key & a set of values for the key
  - It merges these values together to form a smaller set of values
  - Reduce calls are distributed by partitioning the intermediate key space into R pieces using a **partitioning** function (e.g.,  $hash(key) \bmod R$ ). The user specifies the # of partitions (R) and the partitioning function.



# MapReduce

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- **Map**

Grab the relevant data from the source

User function gets called for each chunk of input

- **Reduce**

Aggregate the results

User function gets called for each unique key

# MapReduce: what happens in between?

- **Map**

- Grab the relevant data from the source (parse into key, value)
- Write it to an intermediate file

- **Partition**

- Partitioning: identify which of  $R$  reducers will handle which keys
- Map partitions data to target it to one of  $R$  Reduce workers based on a partitioning function (both  $R$  and partitioning function user defined)

- **Sort**

- Fetch the relevant partition of the output from all mappers
- Sort by keys (different mappers may have output the same key)

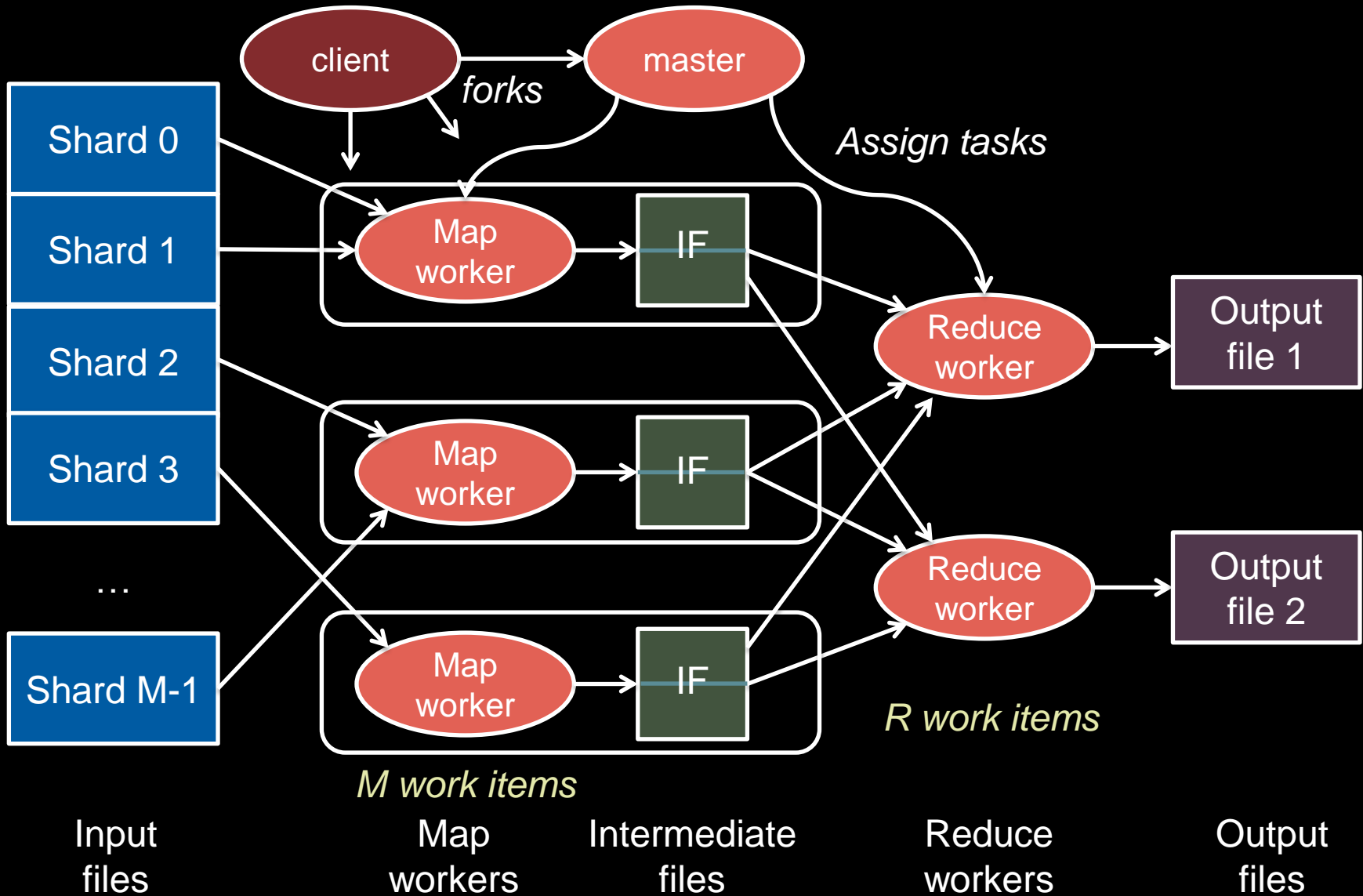
- **Reduce**

- Input is the sorted output of mappers
- Call the user *Reduce* function per key with the list of values for that key to aggregate the results

Map Worker

Reduce Worker

# MapReduce: the complete picture



# Step 1: Split input files into chunks (shards)

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- Break up the input data into  $M$  pieces (typically 64 MB)

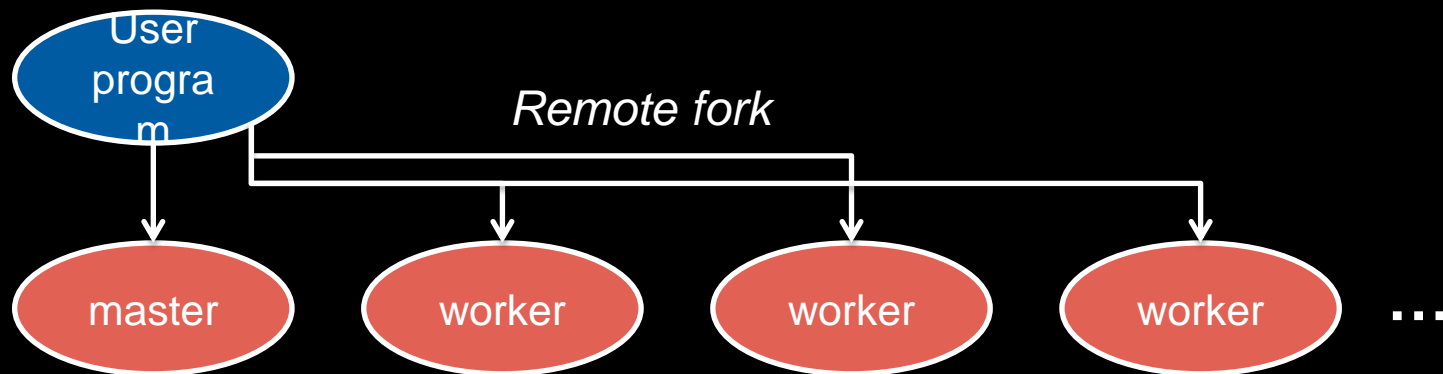


Input files

Divided into  $M$  shards

# Step 2: Fork processes

- Start up many copies of the program on a cluster of machines
  - 1 master: scheduler & coordinator
  - Lots of workers
- Idle workers are assigned either:
  - **map tasks** (each works on a shard) – there are  $M$  map tasks
  - **reduce tasks** (each works on intermediate files) – there are  $R$ 
    - $R = \#$  partitions, defined by the user



# Step 3: Map Task

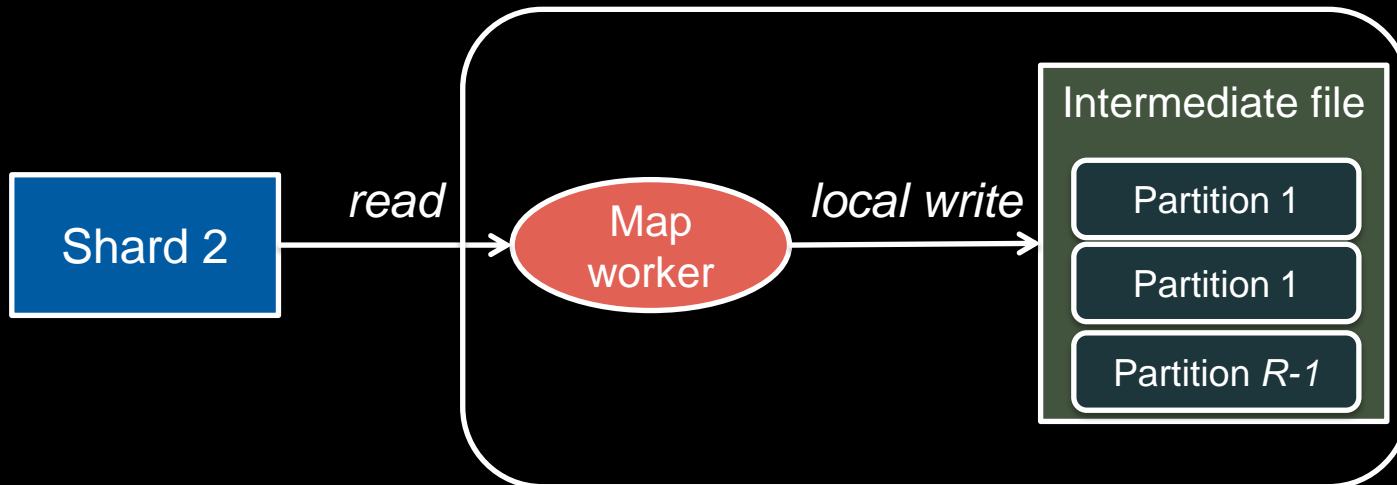
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- Reads contents of the input shard assigned to it
- Parses key/value pairs out of the input data
- Passes each pair to a user-defined *map* function
  - Produces intermediate key/value pairs
  - These are buffered in memory



# Step 4: Create intermediate files

- Intermediate key/value pairs produced by the user's *map* function buffered in memory and are periodically written to the local disk
  - Partitioned into  $R$  regions by a **partitioning function**
- Notifies master when complete
  - Passes locations of intermediate data to the master
  - Master forwards these locations to the reduce worker



# Step 4a. Partitioning

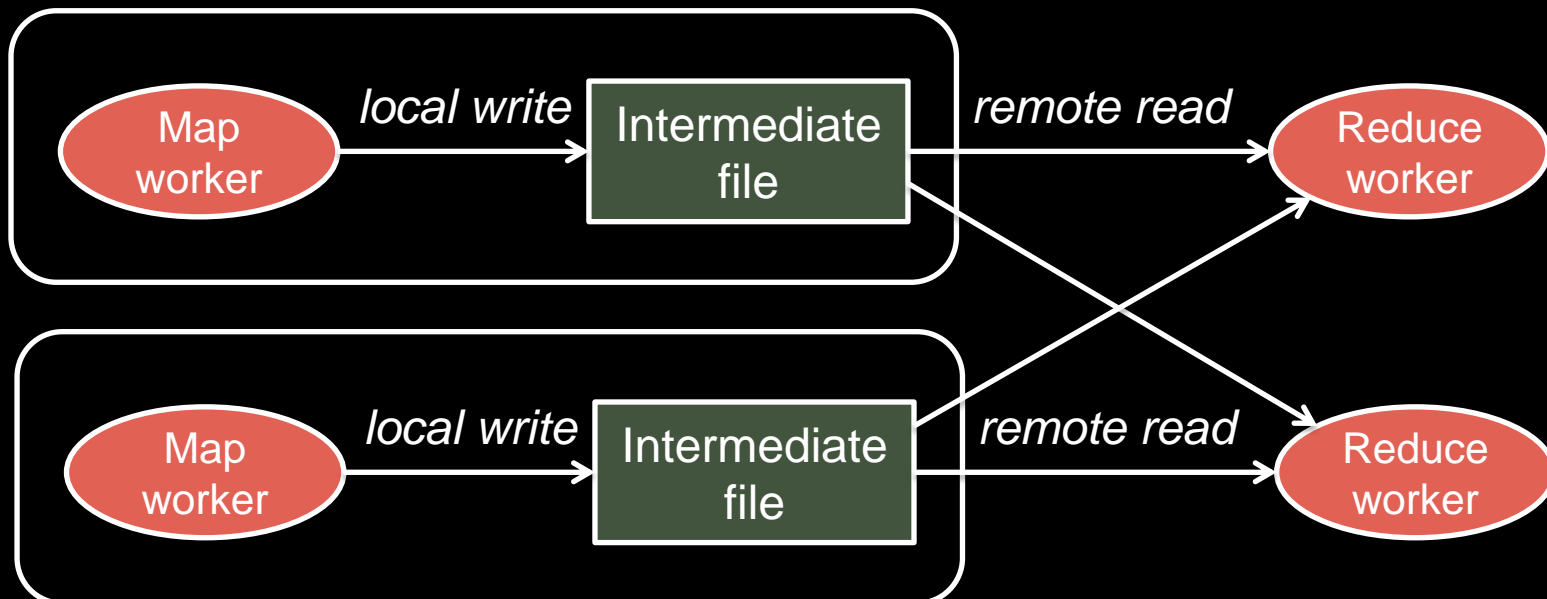
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- Map data will be processed by Reduce workers
  - The user's *Reduce* function will be called once per unique key generated by *Map*.
- This means we will need to sort all the (key, value) data by keys and decide which Reduce worker processes which keys – the Reduce worker will do this
- **Partition function**: decides which of  $R$  reduce workers will work on which key
  - Default function:  $hash(key) \bmod R$
  - Map worker partitions the data by keys
- Each Reduce worker will read their partition from every Map worker



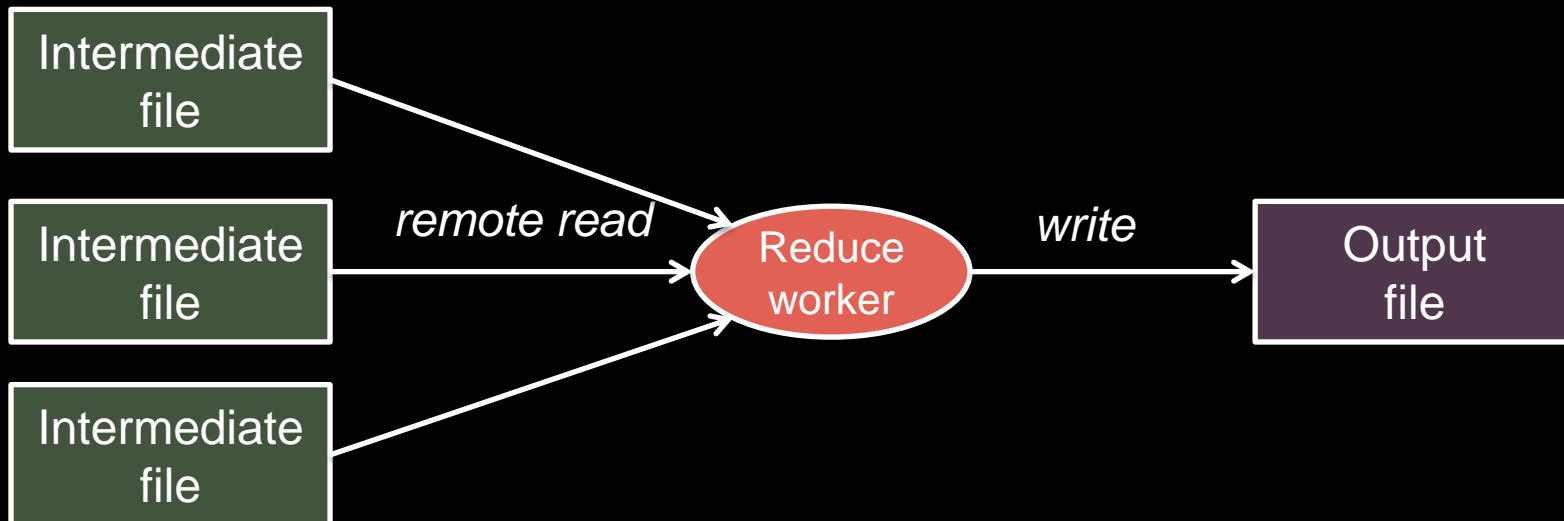
# Step 5: Reduce Task: sorting

- Reduce worker gets notified by the master about the location of intermediate files for its partition
- Uses RPCs to read the data from the local disks of the map workers
- When the *reduce* worker reads intermediate data for its partition
  - It sorts the data by the intermediate keys
  - All occurrences of the same key are grouped together



# Step 6: Reduce Task: *Reduce*

- The sort phase grouped data with a unique intermediate key
- User's *Reduce* function is given the key and the set of intermediate values for that key
  - $\langle \text{key}, (\text{value1}, \text{value2}, \text{value3}, \text{value4}, \dots) \rangle$
- The output of the *Reduce* function is appended to an output file

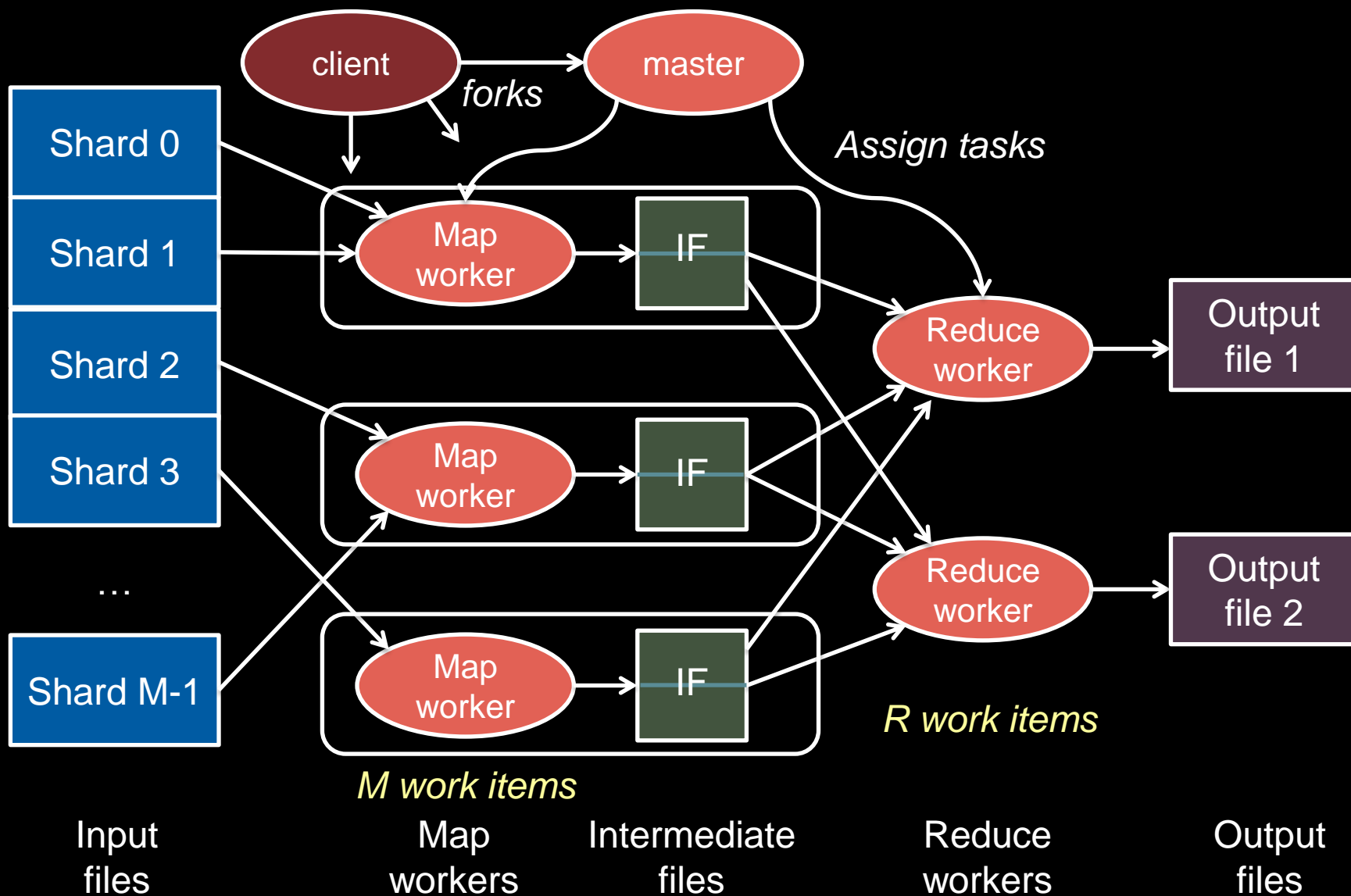


# Step 7: Return to user

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- When all *map* and *reduce* tasks have completed, the master wakes up the user program
- The *MapReduce* call in the user program returns and the program can resume execution.
  - Output of *MapReduce* is available in *R* output files

# MapReduce: the complete picture



# Example

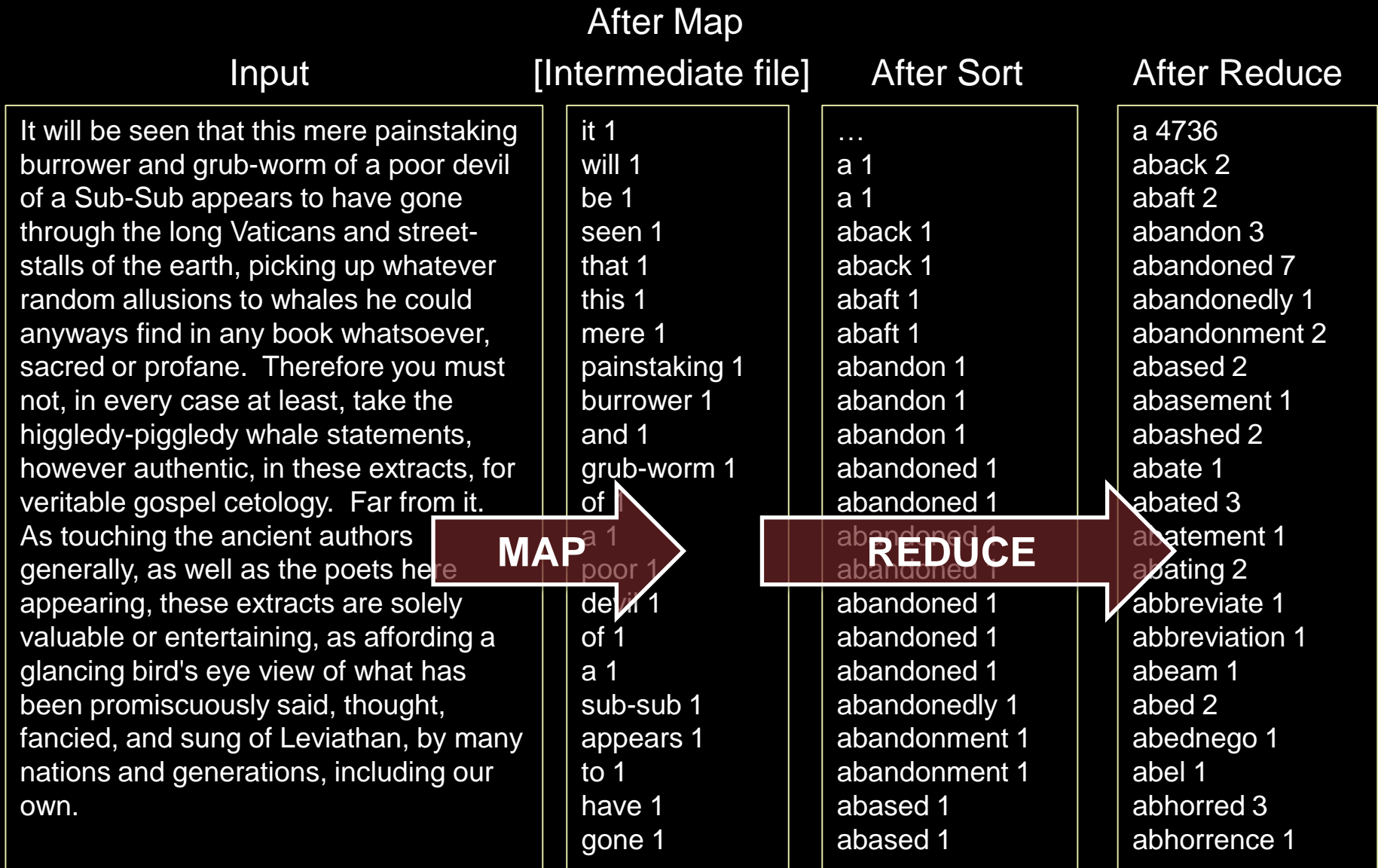
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- Count # occurrences of each word in a collection of documents
- **Map:**
  - Parse data; output each word and a count (1)
- **Reduce:**
  - Sort: sort by keys (words)
  - Reduce: Sum together counts each key (word)

```
map(String key, String value):  
  // key: document name, value: document contents  
  for each word w in value:  
    EmitIntermediate(w, "1");
```

```
reduce(String key, Iterator values):  
  // key: a word; values: a list of counts  
  int result = 0;  
  for each v in values:  
    result += ParseInt(v);  
  Emit(AsString(result));
```

# Example



# Fault tolerance

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- Master pings each worker periodically
  - If no response is received within a certain time, the worker is marked as *failed*
  - *Map* or *reduce* tasks given to this worker are reset back to the initial state and rescheduled for other workers.

# Locality

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- Input and Output files are on GFS (Google File System)
- MapReduce runs on GFS chunkservers
- Master tries to schedule *map* worker on one of the machines that has a copy of the input chunk it needs.



# Other Examples

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- Distributed grep (search for words)
  - *Search for words in lots of documents*
  - Map: emit a line if it matches a given pattern
  - Reduce: just copy the intermediate data to the output

# Other Examples

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- Count URL access frequency
  - *Find the frequency of each URL in web logs*
  - Map: process logs of web page access; output <URL, 1>
  - Reduce: add all values for the same URL

# Other Examples

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- Reverse web-link graph

- *Find where page links come from*
- Map: output  $\langle \text{target}, \text{source} \rangle$  for each link to *target* in a page *source*
- Reduce: concatenate the list of all source URLs associated with a target.

Output  $\langle \text{target}, \text{list}(\text{source}) \rangle$

# Other Examples

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- **Inverted index**

- *Find what documents contain a specific word*
- Map: parse document, emit <word, document-ID> pairs
- Reduce: for each word, sort the corresponding document IDs

Emit a <word, list(document-ID)> pair

The set of all output pairs is an inverted index

# MapReduce Summary

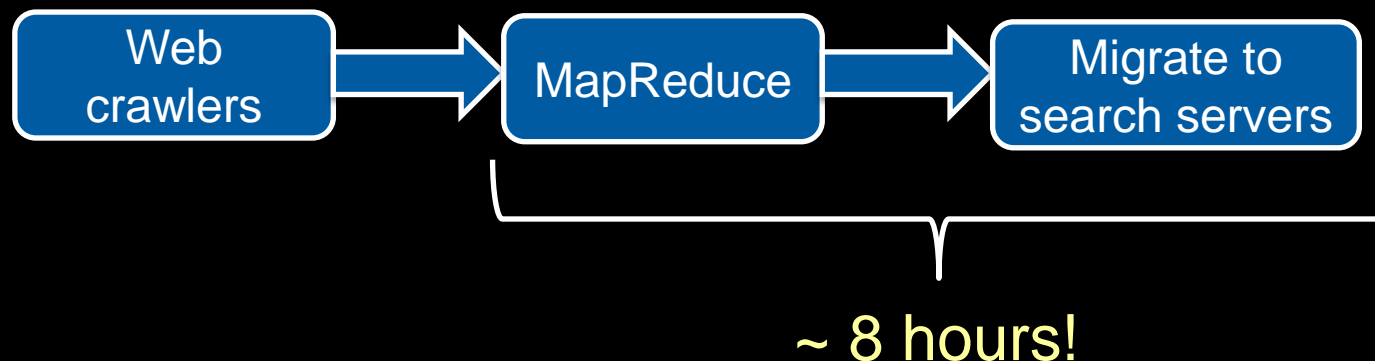
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- Get a lot of data
- **Map**
  - Parse & extract items of interest
- Sort & **partition**
- **Reduce**
  - Aggregate results
- Write to output files

# All is not perfect

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- MapReduce was used to process webpage data collected by Google's crawlers.
  - It would extract the links and metadata needed to search the pages
  - Determine the site's PageRank
- The process took around eight hours.
  - Results were moved to search servers.
  - This was done continuously.



# All is not perfect

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- Web has become more dynamic
  - an 8+ hour delay is a lot for some sites
- Goal: refresh certain pages within seconds
- MapReduce
  - Batch-oriented
  - Not suited for near-real-time processes
  - Cannot start a new phase until the previous has completed
    - Reduce cannot start until all Map workers have completed
  - Suffers from “stragglers” – workers that take too long (or fail)
  - This was done continuously
- MapReduce is still used for many Google services
- Search framework updated in 2009-2010: Caffeine
  - Index updated by making direct changes to data stored in BigTable
  - Data resides in Colossus (GFS2) instead of GFS

# In Practice

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- Most data not simple files
  - B-trees, tables, SQL databases, memory-mapped key-values
- Hardly ever use textual data: slow & hard to parse
  - Most I/O encoded with Protocol Buffers



# More info

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- Good tutorial presentation & examples at:  
<http://research.google.com/pubs/pub36249.html>
- The definitive paper:  
<http://labs.google.com/papers/mapreduce.html>

The End