# **Computer Security**

08r. Exam 1 Review

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# Part 1: Introduction

# Question 1 (B: 6)

#### The CIA Triad is:

- (a) A set of security guidelines established by the U.S. Central Intelligence Agency.
- (b) A collection of techniques hackers use to break into systems.
- (c) Three sets of leaked documents and published on WikiLeaks detailing the CIA's hacking tools.
- (d) A model for classifying topics that need to be addressed in computer security.

#### CIA = Confidentiality, Integrity, Availability

# Question 2 (B: 1)

#### Data confidentiality means that the data:

- (a) Is encrypted.
- (b) Has personally identifiable information removed.
- (c) Cannot be shared without the permission of the owner if it contains personally identifiable information.
- (d) Cannot be accessed by unauthorized parties.

(a) Not always necessary

(b, c) These relate to privacy

# Question 3 (B: 2)

A system with a large *attack surface*:

(a) Offers many ways in which an attacker could try to enter the environment.

- (b) Has a large number of vulnerabilities.
- (c) Is just as likely to be attacked by trusted insiders as well as external attackers.
- (d) Uses multiple forms of defenses to detect and prevent attacks.

Attack surface ≠ number of vulnerabilities

It's a measure of the potential for attack

# Question 4 (B: 3)

#### Advanced Persistent Threats (APT) are most likely to be:

- (a) Small groups of individuals working alone to avoid detection.
- (b) Malicious insiders.
- (c) Intelligence agencies.
- (d) White hat hackers.

#### APTs refer to well-funded, highly determined attackers

# Question 5 (B: 3)

An *opportunistic attack* targets your systems because:

- (a) You are a high-value target.
- (b) Your organization has a malicious insider.
- (c) Attacks from a distance are difficult to trace.
- (d) Your systems may have a vulnerability they are prepared to exploit.

- Opportunistic attacks are the opposite of *targeted attacks*
- The attacker picks you because you are convenient
  - Burglarize a house because the front door is open vs. because the house contains an original Picasso painting

# Question 6 (B: 5)

#### A trusted computing base (TCB) refers to:

- (a) All the components of a system that are critical to its security.
- (b) A computer system that is only available to trusted users.
- (c) Carefully-audited application software that does not interact with non-trusted applications.
- (d) Tamper-resistant computing hardware that the software can trust to run correctly.

TCB = collection of all the hardware, firmware, networks, libraries, programs needed to run an application

### Part 2: Access control

# Question 7 (B: 12)

#### A capability list defines:

(a) The operations that various subjects are allowed to perform on an object.
(b) The system calls that a process can call when it is running with root privileges.
(c) The operations that a subject is allowed to perform on various objects.
(d) The full set of system calls that a process is allowed to invoke.

- Access control list (ACL): associated with objects
  - Defines access rights for various subjects
- Capability list: associated with subjects
  - Defines access rights for various objects

# Question 8 (B: 7)

Unlike access control lists, POSIX (e.g., UNIX, Linux, FreeBSD) permissions: (a) Enumerate a list of users who can access an object.

- (b) Identify a list of objects along with access permissions for those objects.
- (c) Use a fixed amount of space per file to store access permissions.

(d) Allow an administrator to manage group access to objects.

- POSIX file permissions are a restricted form of an access control list
- Each object (file) contains only a set of three subjects: user, group, other

### Question 9 (B: 8)

Unlike discretionary access control (DAC), mandatory access control (MAC):

(a) Requires the kernel to check access rights for an object before opening it.

- (b) Is configured by administrators, not users.
- (c) Organizes users into roles.
- (d) Assigns a confidentiality level to each object.

MAC = <u>mandatory</u> access control

Users cannot override policies set by administrators

- (a) This is done for every form of access control
- (c) This is role-based access control
- (d) This is the Bell-LaPadula model, one form of MAC

# Question 10 (B:9)

Which access model is most directly implemented with an access matrix to manage read/write access rights?

- (a) Bell LaPadula.
- (b) Lattice.
- (c) Biba.
- (d) Type Enforcement.

Type Enforcement is an access control matrix defining access policies of *domains* (groups of users) and *types* (objects)

BLP & Lattice models: read/write access defined based on confidentiality (& compartment) levels

Biba: read/write access defined based on integrity levels

## Question 11 (B: 10)

To which access model would the description, "*because you accessed file A, you now cannot access file B*" apply?

- (a) Lattice.
- (b) Bell LaPadula.
- (c) Chinese wall model.
- (d) Role-based access control (RBAC).

#### The Chinese wall model introduces conflict classes

If you accessed an object that belongs to Group A, you can no longer access objects that belong to Group B if A and B are in a conflict class

# Question 12 (B: 11)

#### The *lattice model*:

- (a) Creates permissions that may change dynamically based on what objects you previously accessed.
- (b) Protects data integrity with a no write up rule.
- (c) Is a form of discretionary access control (DAC).
- (d) Enhances multilevel security.

MAC model with confidentiality hierarchy like Bell-LaPadula

BUT ... compartmentalizes each level to include labels

You need to have a matching label in addition to the allowed level to access data

# Part 3: Injection/hijacking

### Question 13 (B: 19)

#### A **NOP slide** is useful if you:

(a) Are trying to get stack data to overflow onto the heap.

(b) Are using *return-to-libc* techniques instead of code injection.

(c) Want to pad a region of data to prevent the possibility of off-by-one overflows.

(d) Don't know the precise address of your injected code.

- NOP slide = bunch of NOP instructions before the actual attack code
- CPU can branch anywhere into the slide and skip through these instructions until it reaches the useful code

### Question 14 (B: 13)

#### Fuzzing enables:

- (a) Discovery of which input caused a buffer overflow.
- (b) Encryption of pointers to protect them from overflows in the heap.
- (c) Relocation of the starting addresses of the stack, heap, and text (code).
- (d) Run-time buffer overflow checks.

Fuzzing: debugging technique

Supply long chosen data as inputs to cause a buffer overflow

If the program crashes, search for that data in the crash dump

This tells will tell you which input does not check for buffer overflow

### Question 15 (B: 14)

#### Return Oriented Programming (ROP):

- (a) Redirects execution to existing code in the program.
- (b) Forces functions to return prematurely.
- (c) Disables buffer overflow checks.
- (d) Injects executable code onto the stack.

ROP does not inject code but stack frames.

The stack frames contain return addresses to code that is already in the executable (usually a library)

### Question 16 (B: 15)

#### Stack canaries:

- (a) Detect if a program tries to execute code on the stack.
- (b) Ensure that stack data cannot be written outside the stack frame.
- (c) Prevent a function from returning if data on the stack has been corrupted.
- (d) Prevent buffer overflow in stack-allocated variables.

The compiler generates code to:

- 1. Push a random number (a canary) onto the stack
- 2. Check that the canary has not been altered before returning

(a, b, d) They cannot check what happens to data on the stack until the function is ready to return

# Question 17 (B: 16)

What technique is *ineffective* in preventing Return Oriented Programming (ROP) attacks?

- (a) Stack canaries.
- (b) Data execution prevention.
- (c) Address space layout randomization.
- (d) All of the above.

- ROP was designed to bypass no-execute stacks
  - It does not matter if the stack is non-executable because no code is injected into the stack
- Stack canaries detect corruption of stack data
- ASLR make it difficult to inject a valid return address

# Question 18 (B: 17)

*Format string* vulnerabilities arise primarily because:

- (a) User input is used as the format specifier.
- (b) Invalid parameter values are specified.
- (c) More parameters are supplied than the format expects.
- (d) Assumptions are made on the size of the output buffer.

- When users are given the ability to specify the format string, they can insert directives to dump an arbitrary amount of data from the stack or change a value (%n)
- (b, c): parameters are under control of the programmer, not the user
- Possible with *sprintf*, but output is usually to an output stream

### Question 19 (B: 18)

SQL injection attacks cannot be avoided by:

- (a) Not using user input as part of a query or command.
- (b) Escaping all special characters in the input.
- (c) Ensuring the input buffer is sufficiently large to hold the entire query.
- (d) Validating the syntax of the input.

# The problem is user-supplied query data rather than whether it fits into the buffer

### Part 4: Containment

# Question 20 (B: 25)

FreeBSD Jails enhance chroot by::

(a) Allowing multiple applications to share the same jail.

(b) Limiting a jailed application's visible file system to a subtree.

- (c) Restricting the operations allowable to root (admin) within the jail.
- (d) Controlling the system resources (memory, disk) that a jailed process can use.

(a, b) apply to *chroot* as well.

(d) They do not support resource limits

## Question 21 (B: 20)

Linux *namespaces* do not provide the ability to:

- (a) Isolate user IDs.
- (b) Restrict access to system calls.
- (c) Create per-process network stacks.
- (d) Use per-process file system mount points..

**Control groups** restrict system resources

Capabilities restrict what a user can do as root

**Namespaces** provide isolated name spaces

**Chroot jails** are a limited form of namespace that limit what part of the file system you see

# Question 22 (B: 21)

#### Linux **Seccomp-BPF** relies on:

- (a) Using containers to isolate processes.
- (b) Restricting a process' access to only a subtree of the file system space.
- (c) Restrictions on what a process can do when it runs as root.
- (d) Kernel-based restrictions on system calls and file access.

Seccomp-BPF (Berkeley Packet Filter) is a sandboxing kernel extension that allows filtering of system calls and their parameters

- Allow/disallow calls
- Allow/disallow access to specific files, network connections, signals
- (a) They are not containers (which generally provide full isolation)
- (b) This is possible via filtering but is a specific example
- (c) They filter calls regardless of whether a process runs as root or not

### Question 23 (B: 22)

#### The Janus sandbox:

- (a) Uses Linux namespaces for process isolation.
- (b) Provides two levels of sandboxing for greater security: user-level and kernellevel.
- (c) Uses a user-level process to determine if specific system calls should be allowed.
- (d) Validates the operations of a program before it is run to eliminate run-time checks.

Kernel hooks at the system call interface provide callbacks to a user-level process that approves/rejects the call

### Question 24 (B: 23)

#### Linux *capabilities*:

- (a) Restrict what a program can do even if it runs with root privileges.
- (b) Restrict the system calls a program can call even if it is not running with root privileges.
- (c) Allow parts of the file system to be hidden from an application.
- (d) Place limits on the amount of system resources (disk space, network) that a process can consume..

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**Chroot jails** are a limited form of namespace that limit what part of the file system you see

### Question 25 (B: 24)

Unlike containers, *virtual machines* (VMs) can offer applications:

- (a) Separate operating systems.
- (b) Isolated namespaces.
- (c) Simplified packaging of an application and all its dependencies.
- (d) Shared network interfaces.

A VM provides the abstraction of the system hardware (*virtual machine*) and requires an operating system per VM.

### The end