

Computer Security

16r. Pre-exam review

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This covers some highlights of the past four lectures – not all the material

If any of this is really unclear to you, it's an indication that you should spend some time studying the material

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Authentication

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Authentication

- Factors
 1. Something you **have** (key, card, phone, USB dongle)
 2. Something you **know** (password, PIN)
 3. Something you **are** (biometrics)
- Multi-factor authentication
 - Using more than one of these factors
 - E.g., Password + card

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Protocols: Reusable Passwords

Password Authentication Protocol (PAP)

- Classic { *username*, *password* } validation
- **Hashed** passwords
 - Storing $\text{hash}(\text{password})$ ensures that attackers won't see passwords if they get hold of the password file
- **Salted** hashes
 - Adding random text (*salt*) to a password before hashing it guards against dictionary attacks

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Protocols: One-Time Passwords

1. **Sequence-based**: $\text{password} = f(\text{previous password})$
 - Example: S/key authentication
2. **Time-based**: $\text{password} = f(\text{time}, \text{secret})$
 - Example: Time-based One-Time Passwords (TOTP)
3. **Challenge-based**: $f(\text{challenge}, \text{secret})$
 - Example: Challenge-Handshake Authentication Protocol (CHAP)

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U2F: Universal 2nd Factor Authenticator

- Hardware authenticator (usually Bluetooth or USB)
 - Stores public/private keys for each service
- Uses challenge-based authentication
- **Registration with a service** (usually a web site) – initial access
 - Server sends a challenge (*N*)
 - Device generates public/private key pair for the service
 - Sends { *device_id*, *public key*, *N* } signed with its private key
- **Authentication**
 - Server sends a challenge (*N*)
 - Device sends back { *device_id*, *N* } signed with its private key
 - Server can validate using the public key associated with that *device_id*

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Code Signing

Challenge: distribute software and ensure that it is not modified during distribution or on the computer

Solution

- Use **digital signatures**, just like for network messages
- **Publisher:** Hash the software → encrypt the hash with your private key
 - X.509 certificate attached to the application
- **OS:** Hash the software → validate the hash using the publisher's public key
- **Per-page signatures:** sign page-size blocks of software
 - OS can verify a page as it is loaded instead of scanning the entire file ahead of time

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Biometrics

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Biometric Authentication

- Identify a person based on physical or behavioral characteristics
 - Not ownership of keys or knowledge of passwords
- Unlike other forms of authentication
 - Biometrics relies on **statistical pattern recognition**
 - Comparing sampled biometric data with stored biometric data will almost never yield an exact match
 - Software will decide whether the matches are "good enough"
 - False Accept Rate (FAR): false match
 - Non-matching pair of biometric data is *accepted* as a match
 - False Reject Rate (FRR): false non-match
 - Matching pair of biometric data is *rejected* as a match

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Authentication Process

1. **Sensing**
 - Capture the biometric data
2. **Feature extraction**
 - Extract the interesting (unique) parts of the data
3. **Pattern matching**
 - Compare the extracted data with stored samples
4. **Decision**
 - Decide whether the sensed data is close enough to the stored sample


```

    Enrollment: Sensing → Feature extraction → Storage
    Authentication: Sensing → Feature extraction → Matching → Result
  
```

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CAPTCHA

- **Not biometrics** – a technique for software to detect if it's dealing with a human being or a bot
 - Present distorted text (or images) that is difficult for a computer to process but relatively easy for humans
- **Problem:** OCR & computer vision has improved to the point where computers can match human skill
- **NoCAPTCHA RECAPTCHA**
 - No puzzles!
 - Perform "risk analysis"
 - Check IP address of known bots
 - Check Google cookies for legitimate users
 - Track mouse movements for randomness



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Network security

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Data link layer

- **MAC Attacks – CAM overflow**
 - Sniff all data on the local area network*
 - If you send spoofed random source addresses, you will overflow the Ethernet switch's table stored in content-addressable memory (CAM)
 - The switch will then broadcast all traffic onto all ports – *enables sniffing traffic*
- **VLAN Hopping**
 - Sniff all data from connected virtual local area networks*
 - A computer can spoof itself to appear as an ethernet switch with a trunk connection to another switch
 - It will receive traffic for all VLANs (Virtual Local Area Networks) and can see all of it rather than just the traffic for one VLAN

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Data link layer

- **ARP cache poisoning**
 - Redirect IP packets by changing the IP address → MAC address mapping*
 - **Address Resolution Protocol (ARP)**: computer broadcasts a query asking if anyone knows the MAC address corresponding to a given IP address
 - If a malicious host responds with its MAC address, it will receive traffic for that IP address
- **DHCP server spoofing**
 - Configure new devices on the LAN with your choice of DNS address, router address, etc.*
 - Assigns IP address, subnet mask, router address, DNS server address
 - A malicious host can act as a DHCP server and provide bad data for routers or DNS servers to redirect traffic

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Network (IP) & transport (TCP/UDP) layers

No source address authentication – anyone can fake a source address

- **UDP data**– trivial to forge since there is no sequencing
- **TCP data** – harder: need to match sequence numbers
- TCP connection setup
 - **Random starting sequence numbers** make it hard to guess sequence #
 - **SYN flooding attack: Denial of Service**
 - Send TCP connection requests (SYN packets) with an unreachable source address
 - Receiver will allocate resources for the connection
 - Eventually will not be able to accept any more connections
 - Defense: **SYN cookies: minimize SYN flooding problem**
 - Do not allocate resources until the handshake is complete
 - Server computes the SYN-ACK sequence number by
 - `hash(src_addr, dest_addr, src_port, dest_port, SECRET)`
 - SECRET is just a random number that the server picked

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Routing Protocols & DNS

- **BGP** (Border Gateway Protocol)
 - Used by IP networks (autonomous systems) to share **routing** information
 - Uses a TCP connection between routers
 - Route announcements are not authenticated
 - Attacks**
 - **Fake route announcements** can cause routers throughout the Internet to redirect data to a different place
- **DNS** (Domain Name System)
 - Responsible for converting domain names to IP addresses
 - Attacks**
 - Responses can be intercepted & modified, providing the wrong address for a domain name

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Blockchain & Bitcoin

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The blockchain

- Decentralized list of transactions (**ledger**)
 - **Block** = set of transactions (in Bitcoin, ~10-minute window)
 - **Blockchain**: blocks connected via **hash pointers** into a list of blocks
 - Entire blockchain is replicated on all participating servers
 - **Merkle tree**: binary tree of hash pointers within a block to make it easy to find the desired transaction
- User ID (**address**) = your public key
 - Only you have the private key (which is stored in your **wallet**)
- **Guarding against forgery**
 - Each transaction signed by the owner

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Avoiding double spending

- New transactions are sent to all participants
- Each transaction identifies **inputs** (past transactions where the money comes from)
 - No two transactions cannot use the same inputs
 - This ensures there is no **double spending**
- Each participant checks the blockchain to make sure the transaction is valid
- Valid transactions are added to the block

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Proof of Work

When a block is ready to be added to the chain...

Secure the block with a **Proof of Work**

- Field in the block that is modified so that the **hash(block)** has specific properties (first n bits are 0).
- This takes a huge amount of computation – trying different bit patterns

Finding the Proof of Work is called **mining**

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Proof of Work

- The **first server** that computes the Proof of Work advertises it to other systems
 - Each receiver **validates**: this is efficient – just compute the hash
 - Server that finds this gets rewarded with bitcoins
- When a **majority** of systems approves the Proof of Work
 - The block becomes part of the blockchain (connected via a hash pointer to the previous block)

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Changing the Past

The Proof of Work makes it difficult for a server to change old transactions

- You would need to recompute the Proof of Work for all blocks back to the one you need to modify
- This means creating an alternate blockchain
- If there are competing blockchains
 - The **longest chain** is considered the legitimate one
- **51% attack**
 - To alter past transactions & create a longer chain, you need to own over 50% of the computation power

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Confirming transactions

When do we feel safe about a transaction?

- A transaction is **confirmed** after N number of additional blocks are added to the blockchain
- A confirmation value of N mean that an attacker will need to recompute $N+1$ Proof of Work values to modify the blockchain
 - Computationally not feasible
 - Large values of N are recommended for high-value transactions (typically $N=6$)

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Firewalls & VPNs

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Virtual Private Networks

- Key principle: **Tunneling**
- **IPsec** – popular set of VPN protocols
 - **Authentication Header (AH) protocol**
 - Guarantees integrity & authenticity of IP packets – does not encrypt
 - **Encapsulating Security Payload (ESP)**
 - Adds encryption of the entire payload (encapsulated packet)
- IPsec uses
 - **Authentication**: Digital certificates or pre-shared keys
 - **Key exchange**: Diffie-Hellman
 - **Confidentiality**: Symmetric cryptography
 - **Integrity**: HMAC (hash-based MACs)

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Transport Layer Security (TLS)

Goal: provide an authenticated, encrypted, and tamper-proof connection at the **transport layer** between two hosts that software can use in a manner similar to TCP sockets

- **Authentication**
 - Use public key cryptography & X.509 certificates for authentication
- **Key exchange**
 - Diffie-Hellman keys generated per session (or pre-shared keys)
- **Confidentiality**
 - Use symmetric cryptography to encrypt data
- **Integrity**
 - Include an HMAC with transmitted data to ensure message integrity
- Support different key exchange, encryption, integrity, & authentication protocols
 - negotiate what to use at the start of a session

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Firewalls

Firewall (screening router)	1 st generation packet filter that filters packets between networks. Blocks/accepts traffic based on IP addresses, ports, protocols
Stateful inspection firewall	Like a screening router but also takes into account TCP connection state and information from previous connections (e.g., related ports for TCP)
Application proxy	Gateway between two networks for a specific application. Prevents direct connections to the application from outside the network. Responsible for validating the protocol.
IDS/IPS	Can usually do what a stateful inspection firewall does + examine application-layer data for protocol attacks or malicious content
Host-based firewall	Typically screening router with per-application awareness. Sometimes includes anti-virus software for application-layer signature checking
Host-based IPS	Typically allows real-time blocking of remote hosts performing suspicious operations (port scanning, ssh logins)

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Web Security

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Same-origin policy

- Basic security model for the web
 - A browser permits scripts in one page to access data in a second page **only if** both pages have the same origin
 - **Origin** = { URI scheme, hostname, port number }
- Each frame gets the origin of its URL
 - JavaScript code executes with the **authority of its frame's origin**
 - If **cnn.com** loads JavaScript from **jQuery.com**, the script runs with the authority of **cnn.com**
 - Passive content (CSS files, images) has **no** authority
 - It doesn't (and shouldn't) contain executable code
- Cross-Origin Resource Sharing (CORS)
 - A way for the server to tell a browser to treat multiple origins as the same

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Some browser attacks

Cross-Site Request Forgery (CSRF)	<ul style="list-style-type: none"> Browsers send all relevant cookies to the server with each request – these often contain login information Some services place commands on the URL <i>Attack: An attacker may get you to take some action on a site where you are already authenticated by getting you to click on a link</i>
Cross-Site Scripting (XSS)	<ul style="list-style-type: none"> Code injection attack, malicious scripts are added as part of user input and later presented back to a user. Reflected XSS, attacker creates a malicious link. User clicks on it and the response goes back to the user's browser with the malicious script in it. Persistent XSS, attacker adds malicious JavaScript where it will be stored on a server and presented to other users (e.g., forum comments)
Clickjacking	<ul style="list-style-type: none"> Attacker overlays an image to trick a user to clicking a button or link – the user clicks on something different than they think they're clicking on

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Mobile Device Security

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Android Security

- App isolation**
 - Linux user IDs are used as app IDs; each app has its own Linux UID
 - Java apps run in a Dalvik virtual machine
- Mandatory code signing**
 - Can be self-signed or signed by a third party – Android does not validate CAs
- App communication**
 - Apps communicate with **intents**; messages that contain an action & data sent to some other component
 - This is the way apps request services from system services or other apps
- Permissions**
 - Apps must request permission to access resources at install time
 - OS maintains a whitelist of what an app is allowed to access
- File system encryption**

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iOS Security

- App isolation**
 - App sandbox restricts access to other app's data & resources
- App communication**
 - Inter-app communication only through iOS APIs
- Mandatory code signing**
 - Must be signed using an Apple Developer certificate
- App data protection**
 - Apps can use built-in hardware encryption
- File encryption**
 - Each file is encrypted with a unique key

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Hardware protection

- ARM TrustZone: two "worlds"
 - Non-secure world**
 - Cannot access secure resources directly
 - Main OS and apps run in the non-secure (non-trusted) world
 - Secure world**
 - Cryptographic functions & key storage
- Each world has its own OS & applications**
 - Secure key management & key generation
 - Secure boot, digital rights management, secure payment
- Apple Secure Enclave**: Apple's customized TrustZone-like solution
 - Dedicated co-processor for the secure world
 - All cryptographic functions are handled in the secure enclave (secure world)

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The end

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