

Some issues in distributed systems design

- 1. You might need a coordinator
 - But implementing leader election algorithms could be a pain
- 2. You might need to implement locking
 - Implementing distributed mutual exclusion is a pain
- 3. Your collection of programs may need to read configuration info
 - The information may change dynamically (e.g., a group membership list)

Running a central service for locking and object storing is super convenient

- But is not fault-tolerant
- Implementing the service as a replicated state machine will fix that but is a pain
 - Elections, consensus algorithm, RPCs, request redirection, recovery

Google Chubby

Designed as a distributed lock service + simple fault-tolerant file system

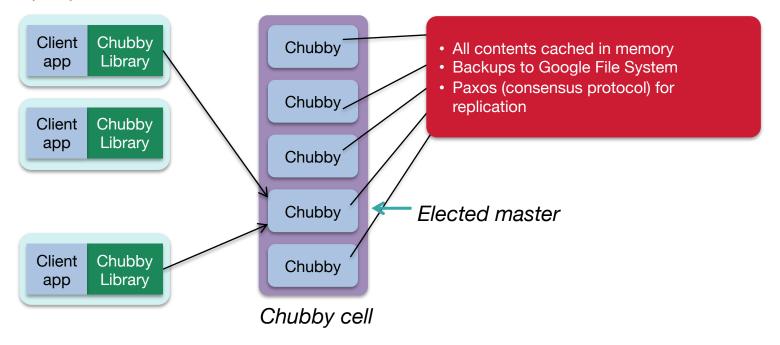
- Interfaces
 - File (object) access
 - Event notification
 - File locking
- Chubby is used to:
 - Manage coarse-grained, long-term locks (hours or days, not < sec)
 - get/release/check lock identified with a name
 - Store small amounts of data associated with a name
 - E.g., system configuration info, identification of primary coordinators
 - Elect masters

Design priority: availability rather than performance

Chubby Deployment

Client library + a Chubby cell (5 replica servers)

Primary/backup replication



Chubby Master

- Chubby has <u>at most</u> one master
 - All requests from the client go to the master

- All other nodes (replicas) must agree on who the master is
 - Paxos consensus protocol used to elect a master
 - Master gets a lease time
 - Re-run master selection after lease time expires to extend the lease
 ...or if the master fails

Simple User-level API for Chubby

- User-level RPC interface
 - Not implemented as an OS-level file system
 - Programs must access Chubby via an API

- Look up Chubby nodes via DNS*
- Ask any Chubby node for the address of the master node
- File system interface (names, content, and locks)

*Domain Name System: the system used to look up internet domain names

Chubby: File System Interface

- /ls/cell/rest/of/name
 - /ls: lock service (common to all Chubby names)
 - cell: resolved to a set of servers in a Chubby cell via DNS lookup
 - /rest/of/name: interpreted within the cell
- Each file has
 - Name
 - Data
 - Access control list
 - Lock
 - No modification, access times
 - No seek or partial reads/writes; no symbolic links; no moves

Think of Chubby as an object store with a hierarchy of names rather than a file system – you cannot read/write partial contents

Chubby: API

open()	Set mode: read, write & lock, change ACL event list, lock-delay, create	,	
close()	Close access to an object		
GetContentsAndStat()	Read file contents & metadata		
SetContents(), SetACL()	Write file contents or ACL		
Delete()	Delete an object		
Acquire(), TryAcquire(), Release()	Lock operations		
GetSequencer()	Sequence # for a lock	Sequence #s are used to detect	
SetSequencer()	Associate a sequencer with a file handle	messages with expired locks. File operations with old sequence #s will be rejected	
CheckSequencer()	Check if sequencer is valid		

Chubby: Files as Locks

- Every file & directory can act as a reader-writer lock
 - Either one client can hold an exclusive (writer) lock
 - Or multiple clients can hold reader locks
- Locks are advisory
- If a client releases a lock, the lock is immediately available
- A client sends a KeepAlive RPC during a Chubby session to say it is still alive
 - If a client fails, the lock will be unavailable for a lock-delay period (typically 1 minute)

Using Locks for Leader Election

Using Chubby locks makes leader election easy

How: a participant tries to acquire a lock

- If it gets the lock, then it's the master for whatever service it's providing!
- No need for user systems to participate in a consensus protocol
 ... the programmer doesn't need to figure out consensus (e.g., Paxos or Raft)
- Chubby provides the needed fault tolerance
- Using locks & sequence numbers to discard old requests
 - Participant gets a lock and requests the lock sequence count from Chubby
 - Share it with other participants.
 - In each RPC to a service, send the sequence count to the server
 - During request processing, a server will reject old (delayed) packets

```
if (sequence_count < current_sequence_count)
    reject request // it must be from a delayed packet</pre>
```

Events

Clients may subscribe to events:

- File content modifications
- Child node added/removed/modified
- Chubby master failed over
- File handle & its lock became invalid
- Lock acquired
- Conflicting lock request from another client

Chubby client caching & master replication

At the client

- Data cached in memory by chubby clients
 - Cache is maintained by a Chubby lease, which can be revoked, invalidating the cache
- All clients write through to the Chubby master

At the master

- Writes are propagated via Paxos consensus to all Chubby replicas
 - Data updated in total order replicas remain synchronized
 - The master replies to a client after the writes reach a majority of replicas
 - Reads can be acknowledged immediately without consulting the replicas.
- Cache invalidations
 - Master keeps a list of what each client may be caching
 - Invalidations are sent by the master and are acknowledged by the client
 - File is then cacheable again
- Chubby database is backed up to GFS every few hours

Chubby – Apache ZooKeeper™

- Chubby is an internal Google service
- Apache ZooKeeper is an open-source centralized service for locking and storing shared services
- Built based on the Chubby paper
 - Uses Zab instead of Paxos for consensus
 (another protocol roughly similar to Paxos but with a focus on log replication, like Raft)
 - Replicated servers
 - Shared hierarchical namespace
 - Stored in memory and organized into files and directories
 - Locking
 - Event notification



Apache ZooKeeper™

- A bit from Chubby but inspired by Chubby & similar in many ways
 - Billed as a "process coordinator" rather than a "lock service"
- Adds "watch mechanism"
 - One-time event trigger to notify a client when a file is updated
- Chubby supports client-side caching with invalidation by the master
 - Zookeeper relies upon clients polling the master or using the watch mechanism to check for updates
 - Zookeeper doesn't guarantee strong consistency:
 - Suppose a client changes the contents of a file and gets an acknowledgment
 - That doesn't mean that all other clients have received notifications of the change even if they are watching the file ⇒ clients may have different ideas of the contents
 - This has been fixed by Netflix with Curator
 - JVM client library for ZooKeeper
 - Adds client caching, strong consistency, and simplified APIs for common tasks

The End